THE PROGRAMMING LANGUAGE PARVA (Level 1)

P.D. Terry, Computer Science Department, Rhodes University

27 June 2003

1 Introduction

Parva is a minimal toy programming language that evolved from Clang (Terry, 1997). Its features are deliberately chosen as a mixture of those found in Pascal and C, so as to support introductory courses in compiler development.

This specification of level 1 of Parva (which does not support multiple function declarations) is modelled on the Modula-2 and Oberon specifications (Wirth (1985), Reiser and Wirth 1992)), but is not intended as a programmer's tutorial. It is intentionally kept concise. Its function is to serve as a reference for programmers, implementors, and manual writers. What remains unsaid is mostly left so intentionally, either because it is derivable from stated rules of the language, or because it would require a general commitment in the definition when a general commitment appears unwise.

2 Syntax

A language is an infinite set of sentences, namely the sentences well formed according to its syntax. Each sentence is a finite sequence of symbols from a finite vocabulary. The vocabulary of Parva consists of identifiers, numbers, strings, operators, delimiters, and comments. These are called lexical symbols and are composed of sequences of characters. (Note the distinction between symbols and characters.)

To describe the syntax the variant of extended Backus-Naur formalism called Cocol/R is used. This is described in full detail elsewhere (Terry, 2005). Brackets [and] denote optionality of the enclosed sentential form, and braces { and } denote its repetition (possibly 0 times). Syntactic entities (non-terminal symbols) are denoted by English words expressing their intuitive meaning. Symbols of the language vocabulary (terminal symbols) are denoted by strings enclosed in quote marks (these include words written in lower-case letters, so-called reserved key words).

3 Vocabulary and representation

The representation of symbols in terms of characters is defined using the ASCII set. Symbols are identifiers, numbers, string literals, character literals, operators, delimiters, and comments.

The following lexical rules must be observed. Blanks and line breaks may appear between symbols but must not occur within symbols (except that line breaks are allowed in comments, and blanks are allowed within string and character literals). They are ignored unless they are essential for separating two consecutive symbols. Capital and lower-case letters are considered as being distinct.

Comments may be inserted between any two symbols in a program. They are arbitrary character sequences either opened by the bracket /* and closed by */, or opened by the bracket // and closed at the end of that line. Comments do not affect the meaning of a program.

```
TOKENS
  identifier = letter { letter | digit | "_" } .
  number = digit { digit } .
  stringLit = '"' { stringCh | backslash printable } '"' .
  charLit = "'" ( charCh | backslash printable ) "'" .
```

Identifiers are sequences of letters and digits. The first character must be a letter.

Examples:

```
x scan Parva Get_Symbol firstLetter
```

Numbers are (unsigned) integers. Integers are sequences of digits, and have their usual decimal interpretation. (The implementation uses 32-bit arithmetic.)

String literals are sequences of zero or more characters or escape sequences enclosed in quote marks ("). The number of characters in a string is called the length of the string. Strings can be used only in fairly limited contexts. A string literal may not extend over a line break in the source text.

Character literals are denoted by a single graphic character or a single escape sequence between single quote marks ('). Character literals denote the integer value of the corresponding ASCII character.

Within a string or character literal the following escape sequences denote non-graphical characters

```
\b Backspace (CHR(8))
\t Horizontal tab (CHR(9))
\n Line feed (CHR(10))
\f Form feed (CHR(12))
\r Carriage return (CHR(13))
\" Quotation mark (CHR(34))
\' Apostrophe (CHR(39))
\x (Where x is not b, t, n, r or f) denotes x itself
```

(Within a string \n denotes the sequence Carriage return, Line feed, on Microsoft systems.)

Examples:

```
"Parva will become Magna" "He said\"food!\" and fed a line\n"
```

Operators and delimiters are the special characters, character pairs, or reserved keywords listed below. The reserved words cannot be used in the role of identifiers. Those in parentheses are not currently keywords, but are reserved for use in later extensions.

```
(default) halt
                    bool
                                                       return
          ! =
&&
                                                       (switch)
                    (break)
                                (do)
                                           i f
     {
                    (case)
                                (else)
                                           int
                                                       true
     }
                    (char)
                                false
                                           new
                                                       void
          <
%
    Γ
                    const
                                (for)
                                           null
                                                       while
                    (continue) (goto)
                                           read
                                                       write
```

4 Programs

A program is a collection of declarations of constants and variables (whose values are said to constitute the program state), and a sequence of statements whose purpose is to alter the program state.

```
Parva = "void" identifier "(" ")" Block .
Block = "{" { Statement } "}" .
```

5 Declarations and scope rules

Every identifier occurring in a program must be introduced by a declaration. Declarations also serve to specify certain permanent properties of an entity, such as its type, and whether it is a constant, a variable, or an array.

The identifier is then used to refer to the associated entity. This is possible only in those parts of a program that are within the scope of the declaration. No identifier may denote more than one entity within a given scope. The scope extends textually from the point of the declaration to the end of the block (see 10.1) in which the declaration has been made and to which the entity is said to be local.

```
Declaration = ConstDeclarations | VarDeclarations .
```

6 Type

In general, a type determines the set of values which variables and constants declared to be of that type may assume, and the operators that are applicable to such variables and constants.

```
Type = BasicType [ "[]" ] .
BasicType = "int" | "bool" .
```

Level 1 of Parva supports only two basic data types, representing the set of Boolean values {false, true} and the set of integer values {min ... max}, where min and max are implementation defined (the implementation uses 32-bit arithmetic {-2147483648 ... 2147483647}).

A type that incorporates the [] token is said to be an *array reference type*. Arrays are constructed as components of expressions by using the keyword new (see section 9.3).

Two entities are said to be *type-compatible* (for the purposes of assignment or comparison) if they are of the same type. In general the type of a value defined by an *Expression* (see section 9) is computed from its operators and the types of its operands; the type of a variable or constant is determined from the type used in its declaration. However, the value denoted by null is a value of all array reference types.

7 Constant declarations

A constant declaration permanently associates an identifier with a constant value.

```
ConstDeclarations = "const" OneConst { "," OneConst } ";" .
OneConst = identifier "=" Constant .
Constant = number | charLit | "true" | "false" | "null" .
```

The type of the identifier is determined from the context of the declaration. true and false are the Boolean constants. Numbers and character literals are deemed to define constants of integer type.

Examples:

```
const
  max = 100;
  YES = true;
  CapitalA = 'A';
```

8 Variable declarations

Variables are those data items whose values may be changed by execution of the program. Variable declarations serve to introduce variables and to associate them with identifiers that must be unique within their given scope. They also serve to associate a fixed data type with each variable so introduced.

```
VarDeclarations = Type OneVar { "," OneVar } ";" .
OneVar = identifier [ "=" Expression ] .
Type = BasicType [ "[]" ] .
BasicType = "int" | "bool" .
```

Each variable identifier denotes either a simple scalar variable of a specified *BasicType* (integer or Boolean), or a reference to an array structure of scalar elements that all have the *BasicType* specified in their declaration. Arrays are constructed as components of expressions by using the keyword new (see section 9.3).

Variables whose identifiers appear in the same list are all of the type specified in the declaration.

When the statement sequence that is defined by a *Block* is activated, all variables are deemed to have initially undefined values, except for those variables that have been assigned the values of expressions within the variable declaration sequence where they were declared.

If a variable is assigned the value of an *Expression* the type of the variable must be compatible with the type of the *Expression*.

Examples:

```
int i, j = 8;
bool startedDating, suitableSpouse;
int[] list = new int[12], list2 = null;
```

Variables used in future examples are assumed to have been declared as indicated above.

9 Expressions

Expressions are constructs denoting rules of computation whereby constants and current values of variables are combined to derive other values by the application of operators. Expressions consist of operands and operators. Parentheses may be used to express specific associations of operators and operands where the normal rules of precedence are unsuitable on their own.

```
Expression = AddExp [ RelOp AddExp ] .

AddExp = [ "+" | "-" ] Term { AddOp Term } .

Term = Factor { MulOp Factor } .

Factor = Designator | Constant | ArrayValue | "!" Factor | "(" Expression ")".

ArrayValue = "new" BasicType "[" Expression "]" .

Designator = identifier [ "[" Expression "]" .

Constant = number | charLit | "true" | "false" | "null" .

AddOp = "+" | "-" | "||" .

MulOp = "*" | "/" | "%" | "&&" .

RelOp = "==" | "!=" | "<=" | ">" | ">=" .
```

9.1 Operands

With the exception of numbers and character literals, many operands in expressions are denoted by designators. A *Designator* consists of an identifier referring to the constant or variable to be designated. This identifier may be followed by a selector if the identifier denotes an array.

```
Designator = identifier [ "[" Expression "]" ] .
```

If A designates an array, then A[E] denotes that element of A whose index is the current value of the expression E. E must be of integer or character type, and the value of E must lie within the range of possible values for the index of A. This range is from 0 ... ArraySize - 1, as specified when A was created.

If the designated entity is a variable or an array element, then, when used as an operand, the designator refers to the variable's current value, and the type of the operand is the type of that variable or array element. An operand that is a literal constant denotes a value that is the value of that constant, and the type of the operand is the type of that literal constant.

An operand may also be designated by the application of the new operator (see section 9.3).

9.2 Operators

The syntax of expressions distinguishes between four classes of operators with different precedences (binding

strengths). The operator ! has the highest precedence, followed by multiplication operators, addition operators, and then relational operators. Operators of the same precedence associate from left to right. Thus, for example, x - y - z * w stands for (x - y) - (z * w).

The available operators are listed in the following tables.

9.2.1 Logical operators

symbol	result
	logical disjunction (Boolean OR) logical conjunction (Boolean AND)
!	negation

These operators apply only when both operands are of the Boolean type, and yield a Boolean result.

It is desirable that the implementation should enforce short-circuit semantics such that

```
p | | q stands for "if p then true, else q"
p && q stands for "if p then q, else false"
! p stands for "not p"
```

However, this has been left as an exercise.

9.2.2 Arithmetic operators

symbol	result
+	sum
_	difference
*	product
/	quotient
%	modulus

These operators apply only to operands of integer type, and yield a result of integer type. When used as operators with a single operand, - denotes sign inversion and + denotes the identity operation.

The operator / produces a result that is truncated towards zero. Implementation of the operator % has been left as an exercise: it is required to produce a result so that the following relation holds for any dividend x and divisor y:

$$x = (x / y) * y + (x % y)$$

Examples

9.2.3 Relational operators

symbol	relation
==	equal
! =	unequal
<	less
<=	less or equal
>	greater
>=	greater or equal

These operators yield a result of Boolean type. The ordering operators <, <=, > and >= apply to operands of the integer type. The relational operators == and != also apply when both operands are of the Boolean type or of the same array reference type (when the comprison is between the values of the references, and *not* between the individual elements of the arrays referred to). The operands of a relational operator may be evaluated in any convenient order.

Examples of expressions:

9.3 Array instantiation

Operands in expressions may have values that are obtained when a new instance of an array is created.

```
ArrayValue = "new" BasicType "[" Expression "]" .
BasicType = "int" | "bool" .
```

An array value allows for the creation of a new instance of an array, whose elements are all of the type specified by *BasicType*. The number of elements is specified by the value of the *Expression*, which must yield a positive value of integer type. The value of the operand is the reference to the newly created array. Typically this is assigned to a variable declared to be of a type specified as BasicType[], whereafter the elements of the array may be accessed by use of appropriate *Designators*.

Examples:

```
int[] list = new int[100];
bool[] sieve = new bool[Max + 1];
```

10 Statements

Statements denote actions or the declaration and possible initialization of variables and constants. Apart from declarative statements, a distinction is drawn between elementary and structured action statements.

Elementary statements are not composed of any parts that are themselves statements. They are the assignment, the input statement, the output statement, the return statement, the halt statement and the empty statement.

Structured statements are composed of parts that are themselves statements. They are used to express sequencing, as well as conditional, selective, and repetitive execution.

A statement may assume the form of a *Block*, and may thus incorporate declarations of identifiers, whose scope is local to that block.

```
Statement = Block | ConstDeclarations | VarDeclarations | Assignment
| IfStatement | WhileStatement | ReturnStatement | HaltStatement
| ReadStatement | WriteStatement
| EmptyStatement .
```

10.1 Blocks (statement sequences)

A *Block* statement may be used to group several statements and declarations into one indivisible unit. Block statements are frequently used as parts of structured statements.

```
Block = "{" { Statement } "}" .
```

Statement sequences within a *Block* statement denote the sequence of actions and declarations specified by the component statements, executed in the order specified by the sequence.

10.2 Assignments

An Assignment denotes the replacement of the value of a variable with a new value, and results in a change to the state of a program.

```
Assignment = Variable "=" Expression ";" .
Variable = Designator .
```

If the *Expression* is present, the assignment serves to replace the current value of the designated *Variable* by the value specified by the *Expression*. The assignment operator is written as "=" and pronounced as "becomes". The types of the *Expression* and of the *Variable* must be compatible, and the *Variable* must designate a scalar variable, an array element or an array reference.

Examples:

```
i = 0;
suitableSpouse = age < 30;
list = new int[max];
list[i] = list[i] + 1;
```

10.3 If statements

An *IfStatement* specifies the conditional execution of a guarded statement.

```
IfStatement = "if" "(" Condition ")" Statement . Condition = Expression .
```

The Expression (guard) must yield a Boolean result. If the guard evaluates to true the guarded Statement is executed.

(There is no else clause in level 1 of Parva. Implementation of this is intended as an exercise.)

Examples:

```
if (i > max) { write("limit exceeded\n"); return; } if (j % 3 = 0) write("divisible by 3");
```

10.4 While statements

A WhileStatement is one form of specifying the repetition of an associated statement.

```
WhileStatement = "while" "(" Condition ")" Statement .
Condition = Expression .
```

The *Expression* that is the *Condition* must yield a Boolean result. If this expression yields true, the associated *Statement* is executed. Evaluation of the condition is performed once at the start of each complete iteration, and the test and statement execution are repeated as long as the *Condition* yields true.

Example:

```
while (j > 0) { j := j / 2; i := i + 1; }
```

10.5 Input statements

```
ReadStatement = "read" "(" ReadElement { "," ReadElement } ")" ";" .
ReadElement = string | Variable .
```

A *ReadStatement* specifies a list of variables that are to be assigned new values from an input source external to the program. Each *Variable* must designate a scalar variable, or an element of an array. A *string* appearing in a *ReadStatement* serves to act as a prompt to the user of the program.

Example:

```
read ("Supply Age", age, "\nSupply Mass", mass);
```

10.6 Output statements

```
WriteStatement = "write" "(" WriteElement { "," WriteElement } ";" .
WriteElement = string | Expression .
```

A WriteStatement specifies a list of strings and expressions whose values are to be computed and then transferred to an output sink external to the program.

Example:

```
write("The result is ", i + j, "\n");
```

10.7 Return statements

A *ReturnStatement* causes immediate termination of the program execution, returning control to the invoking environment.

```
ReturnStatement = "return" ";" .
```

10.8 Halt statements

A *HaltStatement* causes immediate termination of the program execution, returning control to the invoking environment.

```
HaltStatement = "halt" ";" .
```

In level 1 of Parva the *HaltStatement* and the *ReturnStatement* are equivalent. The distinction only becomes apparent in level 2 of Parva.

10.9 Empty statements

An EmptyStatement causes no change to the program state.

```
EmptyStatement = ";" .
```

The *EmptyStatement* is included in order to relax punctuation rules in statement sequences.

11 Complete example

```
void main () {
  const votingAge = 18;
  int age, eligible = 0, total = 0;
  bool allEligible = true;
  int[] voters = new int[100];
  read(age);
  while (age > 0) {
    bool canVote = age > votingAge;
    allEligible = allEligible && canVote;
    if (canVote) {
      voters[eligible] = age;
      eligible = eligible + 1;
      total = total + voters[eligible - 1];
    }
    read(age);
}
  write(eligible, " voters. Average age is ", total / eligible, "\n");
  if (allEligible) write("Everyone was above voting age");
}
```

12 Missing statements

There are no do, goto, for, break, switch or continue statements in this release of the language. Their implementation is left as a series of exercises!

13 Bibliography

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